

# A pattern matcher for miniKanren *or*

## How to get into trouble with CPS macros

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# Introduction

- Writing a pattern matcher for miniKanren
  - Automatically bind variables from pattern
- A CPS macro implementation
  - Delayed binding and conditional expansion
  - Auxiliary keyword comparison
- Two fixes for the CPS macro
  - `bound-identifier=?` with **`syntax-case`**
  - Eager binding

# miniKanren

- A relational programming language
- Embedded in Scheme
- Similar to Prolog, different approach
  - Interleaved search
  - Lexical scoping of variables

# CPS Macros

- Stay within **syntax-rules**
  - Simpler semantics
  - Provides powerful macro facility
- Control expansion order
- Allows for conditional expansion
  - Some caution may be required

# Pattern matcher goals

- Simplify miniKanren relations
  - Provide simple pattern matching
  - Automatically create variables in pattern
- Generate relations with good performance
  - “Fail fast” philosophy, tight variable scope
  - Avoid creating unnecessary binding forms

# An Example: append

```
append([], Y, Y) .  
append([A|D], Y2, [A|R])  
    :- append(D, Y2, R) .
```

# An Example: append

```
(define append
  (λ (x y z)
    (conde
      ((≡ `() x) (≡ y z))
      ((exist (a d r)
        (≡ `(,a . ,d) x)
        (≡ `(,a . ,r) z)
        (append d y r))))))
```

# An Example: append

```
(define append
  ( $\lambda^e$  (x y z)
    (( () ___ , y) )
    (( ( , a . , d) ___ ( , a . , r) )
      (append d y r) ) ) )
```



# An Example: append

```
(define append
  ( $\lambda^e$  (x y z)
    (( () ___ , y) )
    (( ( , a . , d) ___ ( , a . , r) )
      (append d y r) ) ) )
```

---

```
append ( [] , Y , Y ) .
```

```
append ( [A | D] , Y2 , [A | R] )
```

```
:- append ( D , Y2 , R ) .
```

# Original Algorithm

- $\lambda^e$  and  $\text{match}^e$
- `handle-clauses`
- `handle-pattern`
- `build-cons`
- `build-goal`
- `build-var`
- `build-clauses-part`
- `build-clause`
- `make-clauses-cont`
- `make-pattern-cont`
- `case-id`

# Original algorithm

- CPS macro controls expansion order
- Binding delayed until end of expansion
- Unbound identifiers compared
  - free, symbolically-equal identifiers appear equal
  - Otherwise equal identifiers introduced in different lexical scopes left unbound

# Original algorithm


```
(define append
  (λe (x y z)
    ( ( ( ) ___ , y ) )
    ( ( ( , a . , d ) ___ ( , a . , r ) )
      (append d y r) ) ) )
```

# Original algorithm

```
(define append  
  ( $\lambda^e$  (x) y z)  
    ( ( ( ) _____ , y) )  
    ( ( ( , a . , d) _____ ( , a . , r) )  
      (append d y r) ) ) )
```

```
( (ex a d) (= (cons a d) x) )
```

# Original algorithm

```
(define append  
  ( $\lambda^e$  (x y z)  
    (( ()        , y) )  
    (( ( , a . , d)  ( , a . , r) )  
      (append d y r) ) ) )
```

```
( (ex a d) (= (cons a d) x) )
```

# Original algorithm

```
(define append
  (λe (x y z)
    (( () ___ , y) )
    (( ( , a . , d) ___ ( , a . , r) )
      (append d y r) ) ) )
```

```
((ex a d) (= (cons a d) x)
 (ex r) (= (cons a r) z) )
```

# Original algorithm

```
(define append
  ( $\lambda^e$  (x y z)
    (( () ___ , y) )
    (( ( , a . , d) ___ ( , a . , r) )
      (append d y r) ) ) )

(exist (a d)
  ( $\equiv$  (cons a d) x)
  (exist (r)
    ( $\equiv$  (cons a r) z)
    (append d y r) ) )
```



# Original Algorithm

```
(define append
  (λ (x y z)
    (conde
      ((≡ '() x) (≡ y z))
      ((exist (a d)
        (≡ (cons a d) x)
        (exist (r)
          (≡ (cons a r) z)
          (append d y r)))))))
```

# Original Algorithm

```
(define append
  (λ (x y z)
    (conde
      ((≡ `() x) (≡ y z))
      ((exist (a d r)
        (≡ `(,a . ,d) x)
        (≡ `(,a . ,r) z)
        (append d y r))))))

(define append
  (λ (x y z)
    (conde
      ((≡ '() x) (≡ y z))
      ((exist (a d)
        (≡ (cons a d) x)
        (exist (r)
          (≡ (cons a r) z)
          (append d y r))))))
```

# Exposing our bug

```
(define-syntax break-λe
  (syntax-rules ()
    ( ( _ v)
      (λe (x y) ( ( ( , w . , v) , v) ) ) ) ) )
```

(break-λ<sup>e</sup> z) ⇒

(λ<sup>e</sup> (x y) ( ( ( , w . , z) , z) ) ) ⇒

(λ (x y)

(cond<sup>e</sup>

( (exist (z w)

(== (cons w z) x)

(== z y) ) ) ) )

# Exposing our bug

```
(define-syntax break-λe
  (syntax-rules ()
    ( ( _ v)
      (λe (x) y) ( ( ( , w . , v) , v) ) ) ) )
```

```
(break-λe x) ⇒
(λe (x) y) ( ( ( , w . , x) , x) ) ⇒
(λ (x) y)
  (conde
    ( (exist (x) w)
      (== (cons w x) x)
      (== x y) ) ) )
```

# Exposing our bug

```
(define-syntax break-λe
  (syntax-rules ()
    ( ( _ v)
      (λe (x y) ( ( ( , W) . , v) , v) ) ) ) )
```

(break-λ<sup>e</sup> W) ⇒

(λ<sup>e</sup> (x y) ( ( ( , W) . , W) , W) ) ⇒

(λ (x y)

(cond<sup>e</sup>

( (exist (W)

(== (cons W W) x)

(== W y) ) ) )

# Source of our bug

```
(define-syntax case-id
  (syntax-rules (else)
    (( _ x ((x** ...) act*) ... (else e-act) )
     (letrec-syntax
       ((helper
         (syntax-rules (x else)
           (( _ (else a) ) a)
           (( _ (() a) c . c*) (helper c . c*))
           (( _ (x . z*) a) c . c*) a)
           (( _ (y z* (... ...)) a) c . c*)
              (helper ((z* (... ...)) a) c . c*))))))
    (helper
      ((x** ...) act*) ... (else e-act))))))
```

# Source of our bug

```
(define-syntax case-id
  (syntax-rules (else)
    ((_ x ((x** ...) act*) ... (else e-act))
     (letrec-syntax
       ((helper
         (syntax-rules (x else)
           ((_ (else a)) a)
           ((_ (() a) c . c*) (helper c . c*))
           ((_ (x . z*) a) c . c*) a)
           ((_ (y z* (... ...)) a) c . c*)
              (helper ((z* (... ...)) a) c . c*))))))
    (helper
      ((x** ...) act*) ... (else e-act))))))
```

# Source of our bug

```
(define-syntax case-id
  (syntax-rules (else)
    ((_ x ((x** ...) act*) ... (else e-act))
     (letrec-syntax
       ((helper
         (syntax-rules (x) else)
           ((_ (else a)) a)
            ((_ (() a) c . c*) (helper c . c*))
            ((_ (x . z*) a) c . c*) a)
            ((_ ((y z* (... ..)) a) c . c*)
              (helper ((z* (... ..)) a) c . c*))))))
    (helper
      ((x** ...) act*) ... (else e-act))))))
```



# Source of our bug

- Auxiliary keyword check is similar to `free-identifier=?`
- Delayed binding leaves variables free
- Free, symbolically-equal identifiers may be conflated

# bound-identifier=? to the rescue

- Identifiers compared as if they were bound
  - Free, symbolically-equal variables introduced in different lexical scopes can be distinguished
- Needs **syntax-case**
  - More than term-rewriting
  - Needs Scheme conditionals to be useful

# case-id redux

```
(define-syntax case-id
  (λ (exp)
    (syntax-case exp (else)
      ((_ x (else e-act)) #'e-act)
      ((_ x ((y x* ...) act) ((x** ...) act*) ...)
       (else e-act))
      ((bound-identifier=? #'x #'y)
       #'act)
      ((_ x ((y x* ...) act) ((x** ...) act*) ...)
       (else e-act))
      #'(case-id x
          ((x* ...) act) ((x** ...) act*) ...)
        (else e-act)))
      ((_ x (() act) ((x** ...) act*) ...)
       (else e-act))
      #'(case-id x
          ((x** ...) act*) ... (else e-act))))))
```

# Eager binding to the rescue

- Binds variables as they are encountered
  - Ensures at most one variable free in comparison
  - `syntax-rules` based `case-id` sufficient
- Requires rewrite of code handling identifier
  - Expression built as expansion happens
  - Binding forms can be recognized by expander

# Eager binding to the rescue

- $\lambda^e$  and `matche`
- `handle-clauses`
- `do-clause`
- `do-pattern-opt`
- `do-pattern`
- `handle-pattern`
- `handle-pattern-cont`
- `exist-helper`

# Comparing Output

```
(define append
  (λ (x y z)
    (conde
      ((≡ '() x) (≡ y z))
      ((exist (a d)
        (≡ (cons a d) x)
        (exist (r)
          (≡ (cons a r) z)
          (append d y r))))))
```

# Comparing Output

```
(define append
  ( $\lambda$  (x y z)
    (conde
      ((exist ()) ( $\equiv$  () x) ( $\equiv$  y z)))
      ((exist (a)
        (exist (d)
          (exist ()
            ( $\equiv$  (cons a d) x)
            (exist (r)
              (exist ()
                ( $\equiv$  (cons a r) z)
                (append d y r))))))))))
```

# Summary

- CPS macros powerful, require care
  - Limited to `free-identifier=?`
  - Conditional expansion on identifiers to be bound should be done with caution
- `bound-identifier=?` is a simple fix
- Possible to stay within **`syntax-rules`**
- Larger job to rewrite for eager binding



# Thanks

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  - Lindsey Kuper
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# Questions?

